

On the edge r -dynamic chromatic number of some related graph operations

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Abstract—All graphs in this paper are simple, nontrivial, connected and undirected. By an edge proper k -coloring of a graph G , we mean a map $c : E(G) \rightarrow S$, where $|S| = k$, such that any two adjacent edges receive different colors. An edge r -dynamic k -coloring is a proper k -coloring c of G such that $|c(N(uv))| \geq \min(r, d(u) + d(v) - 2)$ for each edge uv in $V(G)$, where $N(uv)$ is the neighborhood of uv and $c(S) = c(uv) : uv \in S$ for an edge subset S . The edge r -dynamic chromatic number, written as $\lambda_r(G)$, is the minimum k such that G has an edge r -dynamic k -coloring. In this paper, we will determine the edge coloring r -dynamic number of a comb product of some graph, denote by $G \triangleright H$. Comb product of some graph is a graph formed by two graphs G and H , where each edge of graph G is replaced by which one edge of graph H .

Keywords— r -dynamic chromatic number, graph coloring, exponential graphs

INTRODUCTION

Graph theory is part of discrete mathematics which is mostly studied by researchers. This is due to the wide of application in real life. One of the theory developed in graph theory is the coloring. The new extension of graph color is r -dynamic coloring. The r -dynamic chromatic number, introduced by Montgomery [1] and written as $\chi_r(G)$, is the least k such that G has an r -dynamic k -coloring. Note that the 1-dynamic chromatic number of graph is equal to its chromatic number, denoted by $\chi(G)$, and the 2-dynamic chromatic number of graph has been studied under the name a dynamic chromatic number, denoted by $\chi_d(G)$. In [1], he conjectured $\chi_2(G) \leq \chi(G) + 2$ when G is regular, which remains open. Akbari *et.al.* [2] proved Montgomery's conjecture for bipartite regular graphs, as well as Lai, *et.al.* [3] proved $\chi_2(G) \leq \Delta(G) + 1$ for $\Delta(G) \leq 3$ when no component is the 5-cycle.

By a greedy coloring algorithm, Jahanbekama [4] proved that $\chi_r(G) \leq r\Delta(G) + 1$, and equality holds for $\Delta(G) > 2$ if and only if G is r -regular with diameter 2 and girth 5. They improved the bound to $\chi_r(G) \leq \Delta(G) + 2r - 2$ when $\delta(G) > 2r \ln n$ and $\chi_r(G) \leq \Delta(G) + r$ when $\delta(G) > r^2 \ln n$.

The following observation is useful to find the exact values of r -dynamic chromatic number.

Observation 1. Let $\delta(G)$ and $\Delta(G)$ be a minimum and maximum degree of a graph G , respectively. Then the followings hold

- $\chi_r(G) \geq \min\{\Delta(G), r\} + 1$,
- $\chi(G) \leq \chi_2(G) \leq \chi_3(G) \leq \dots \leq \chi_{\Delta(G)}(G)$,
- $\chi_{r+1}(G) \geq \chi_r(G)$ and if $r \geq \Delta(G)$ then $\chi_r(G) = \chi_{\Delta(G)}(G)$.

There are some researches studied this problem, some of them can be found in [5],[6],[7],[8],[9].

THE RESULTS

We are ready to show our main theorems. In this study we used exponential graph, there are two theorems found in this study. These deals with exponential graph $C_n \triangleright P_2$ and $C_n \triangleright C_m$.

Theorem 1. Let G be a comb product denote by $C_n \triangleright P_2$ for $n \geq 3$, the vertex r -dynamic chromatic number is:

$$\chi(C_n \triangleright P_2) = \begin{cases} 3, & n \text{ even} \\ 3, & n \text{ odd} \end{cases}$$

$$\chi_d(C_n \triangleright P_2) = \begin{cases} 3, & n = 3k, k \in \mathbb{N} \\ 4, & n \text{ otherwise} \end{cases}$$

$$\chi_{r \geq 3}(C_n \triangleright P_2) = \begin{cases} 4, & n = 3k, n = 4k, k \in \mathbb{N} \\ 4, & n \text{ otherwise} \\ 5, & n = 5k, k \in \mathbb{N} \end{cases}$$

Proof. A comb product of graph $C_n \triangleright P_2$, for $n \geq 3$, is a connected graph with vertex set $V(C_n \triangleright P_2) = \{x_i, 1 \leq i \leq n\} \cup \{y_i, 1 \leq i \leq n\}$, and edge set $E(C_n \triangleright P_2) = \{x_i x_{i+1}; 1 \leq i \leq n-1\} \cup \{x_i y_i; 1 \leq i \leq n\}$. The order and size of $C_n \triangleright P_2, n \geq 3$ are $|V(C_n \triangleright P_2)| = 2n$ and $|E(C_n \triangleright P_2)| = 2n$. A comb product graph

$C_n \triangleright P_2$ is regular graph of degree 3, thus $C_n \triangleright P_2$,

$\delta(C_n \triangleright P_2) = \Delta(C_n \triangleright P_2) = 3$. By observation 1, $\chi_r(C_n \triangleright P_2) \geq \min\{\Delta(C_n \triangleright P_2), r\} = \min\{3, r\}$. To find the exact value of r -dynamic chromatic number of $C_n \triangleright P_2$, we define three cases, namely $\chi(C_n \triangleright P_2)$, $\chi_2(C_n \triangleright P_2)$ and $\chi_{r \geq 3}(C_n \triangleright P_2)$.

For $r = 1$, the lower bound $\chi(C_n \triangleright P_2) \geq \min\{3, 1\} = 1$, and for $r = 2$, the lower bound $\chi(C_n \triangleright P_2) \geq \min\{3, 2\} = 2$. We will prove that $\chi(C_n \triangleright P_2) \leq 3$ by defining a map $c_1 : V(C_n \triangleright P_2) \rightarrow \{1, 2, \dots, k\}$ for $n \geq 3$, by the following:

$$\begin{aligned} & 12 \dots 12, n \text{ even} \\ & 123 \dots 123, n \text{ odd} \end{aligned}$$

$$c_1(y_1, y_2, \dots, y_n) = \begin{cases} 21 \dots 21, n \text{ even} \\ 212 \dots 212, n \text{ odd} \end{cases}$$

It is easy to see that c_1 gives $\chi(C_n \triangleright P_2) \leq 2$ for n even, but for n odd, we could not avoid to have $\chi(C_n \triangleright P_2) \leq 3$, otherwise there are at least two adjacent vertices assigned the same colors. Thus $\chi(C_n \triangleright P_2) = 2$ for n even and $\chi(C_n \triangleright P_2) \leq 3$, for n odd.

For $\chi_r(C_n \triangleright P_2)$ and $r \geq 3$, the lower bound $\chi_3(C_n \triangleright P_2) \geq \min\{3, 3\} = 3$. We will prove that $\chi_3(C_n \triangleright P_2) \leq 4$ by defining a map $c_3 : V(C_n \triangleright P_2) \rightarrow \{1, 2, \dots, k\}$ for $n \geq 3$, by the following.

$$c_3(x_1, x_2, \dots, x_n) = \begin{cases} 123 \dots 123, n = 3k, \\ 1234 \dots 1234, n = 4k, \\ 12345 \dots 12345, \\ n = 5k, \\ 1231234, n = 7, n \equiv 7 \pmod{3}, \\ 12341234123, n = 11, \end{cases}$$

$$c_3(y_1, y_2, \dots, y_n) = \begin{cases} 444 \dots 444, & n = 3k, \\ 3412 \dots 3412, & n = 4k, \\ 45123 \dots 45123, \\ n = 5k, \\ 3444412, & n = 7, \\ n \equiv 7 \pmod{3}, \\ 44123412344, & n = 11, \end{cases}$$

It is easy to see that c_3 gives $\chi_3(C_n \supseteq P_2) \leq 4$, for $n = 3k, n = 4k, k \in N$, but for $n = 5k$ we are forced to have $\chi_3(C_n \supseteq P_2) = 5$ as well as $\chi_3(C_n \supseteq P_2) = 4$ for n otherwise.

Thus $\chi_3(C_n \supseteq P_2) = 4$, for $n = 4k$, and $\chi_3(C_n \supseteq P_2) = 5$ for n otherwise. By observation 1 $r \geq \Delta(C_n \supseteq P_2) = 3$, it immediately gives $\chi_3(C_n \supseteq P_2) = \chi_r(C_n \supseteq P_2)$ for $n \geq 3$. \square

Theorem 2. Let G be a comb product denote by $C_n \supseteq P_2$ for $n \geq 3$, the edges r -dynamic chromatic number is:

$$\lambda(C_n \supseteq P_2) = \lambda_d(C_n \supseteq P_2) = \lambda_3(C_n \supseteq P_2) = 3$$

$$\lambda_{r \geq 3}(C_n \supseteq P_2) = 4$$

Proof. A comb product of graph $C_n \supseteq P_2$, for $n \geq 3$, is a connected graph with vertex set $V(C_n \supseteq P_2) = \{x_i, 1 \leq i \leq n\} \cup \{y_i, 1 \leq i \leq n\}$, and edge set $E(C_n \supseteq P_2) = \{x_i x_{i+1}; 1 \leq i \leq n-1\} \cup \{x_i y_i; 1 \leq i \leq n\}$. The order and size of $C_n \supseteq P_2, n \geq 3$ are $|V(C_n \supseteq P_2)| = 2n$ and $|E(C_n \supseteq P_2)| = 2n$. A comb product of graph $C_n \supseteq P_2$ is regular graph of degree 3, thus $C_n \supseteq P_2$,

$\delta(C_n \supseteq P_2) = \Delta(C_n \supseteq P_2) = 3$. By observation when $\chi_r(C_n \supseteq P_2) \geq \min\{\Delta(C_n \supseteq P_2), r\} = \min\{3, r\}$. To find the exact value of r -dynamic chromatic number of $C_n \supseteq P_2$, we define three cases, namely $\chi(C_n \supseteq P_2), \chi_2(C_n \supseteq P_2)$ and $\chi_{r \geq 3}(C_n \supseteq P_2)$.

Case 1. For $r \geq 2$, the lower bound $\Delta(G) \leq \chi(G) \leq \Delta(G) + 1$, that $\chi(C_n \supseteq P_2) \geq 3$. Furthermore, to show that $\chi(C_n \supseteq P_2) \leq 3$ with coloring edges $E(C_n \supseteq P_2)$ as in function c_1 . Let $D = \{1, 2, \dots, k\}$ is set of color from k -coloring that c_1 the function by defining a map edges coloring $D, c_1 : E(C_n \supseteq P_2) \rightarrow D$, so mapping each edges to set color D , by the following that:

$$c_1(x_1 y_1, x_2 y_2, \dots, x_n y_n) = \begin{cases} 12 \dots 12, \\ n \text{ even} \\ 12 \dots 12 3, \\ n \text{ odd} \\ 2 33 \dots 33 1, \\ n \text{ odd} \\ 33 \dots 33, \\ n \text{ even} \end{cases}$$

From function coloring c_1 seen that the chromatic number is $\chi(C_n \supseteq P_2) \leq 3$. Because $\chi(C_n \supseteq P_2) \leq 3$ and $\chi(C_n \supseteq P_2) \geq 3$, then $\chi(C_n \supseteq P_2) = 3$, so that $\chi(C_n \supseteq P_2) = \chi_2(C_n \supseteq P_2) = 3$.

Case 2. For $r \geq 3$, the lower bound $\chi_3(G) \geq \chi_2(G)$, that $\chi_3(C_n \supseteq P_2) \geq 4$. We will prove that $\chi_3(C_n \supseteq P_2) = 3$ such as coloring function c_1 that the definition edges coloring r -dynamic not full filled. It is caused by edge $x_1 x_2$, earned $d(u) + d(v) - 2 = 4, |c(N(e))| = 2$ dan $\min\{r, d(u) + d(v) - 2\} = \min\{3, 4\} = 3$, so that $2 \not\geq 3$. So, lower bound is $\chi_3(C_n \supseteq P_2) \geq 4$. Furthermore, to show that $\chi_3(C_n \supseteq P_2) \leq 4$ with coloring edge $E(C_n \supseteq P_2)$ as in function c_2 . Let $D = \{1, 2, \dots, k\}$ is set of color from k -coloring that c_2 the function by defining a map edges coloring $D, c_2 : E(C_n \supseteq P_2) \rightarrow D$, so mapping

each edges to set color D , by the following that:

$$c_2(x_1 x_2, x_2 x_3, \dots, x_n x_1) = \begin{cases} 12 \dots 12, \\ n \text{ even} \\ 12 \dots 12 3, \\ n \text{ odd} \\ 43 \dots 43 4, \\ n \text{ odd} \\ 34 \dots 34, \\ n \text{ even} \end{cases}$$

From function coloring c_2 seen that the chromatic number is $\chi(C_n \supseteq P_2) \leq 4$. Because $\chi(C_n \supseteq P_2) \leq 4$ and $\chi(C_n \supseteq P_2) \geq 4$, then $\chi(C_n \supseteq P_2) = 4$, so that $\chi(C_n \supseteq P_2) = \chi_{r \geq 3}(C_n \supseteq P_2) = 4$. \square

Theorem 3. Let G be a comb product denote by $C_n \supseteq C_m$ for $n \geq 3$ and $m \geq 3$, vertex r -dynamic chromatic number of $C_n \supseteq C_m$ is:

$$\chi(C_n \supseteq C_m) = \begin{cases} 2, & n \text{ even} \\ 3, & n \text{ otherwise} \end{cases}$$

$$\chi_2(C_n \supseteq C_m) = \begin{cases} 3, & n = 3, m = 3 \end{cases}$$

$$\chi_3(C_n \supseteq C_m) = 4, n = 3, m = 3$$

$$\chi_{r \geq 2}(C_n \supseteq C_m) = \begin{cases} 4, & n = 5, m = 5 \\ 5, & n = 3k, m = 3k \end{cases}$$

$$\chi_{r \geq 4}(C_n \supseteq C_m) = \begin{cases} 5, & n = 3k, i \leq i \leq n, \\ & 1 \leq j \leq m - 2 \end{cases}$$

Proof. The comb product of graph denoted by $C_n \supseteq C_m$, is connected graph with vertex set $V(C_n \supseteq C_m) = \{x_i; 1 \leq i \leq n\} \cup \{x_{i,j}; 1 \leq i \leq n; 1 \leq j \leq m-2\}$ and edge set $E(C_n \supseteq C_m) = \{x_n x_1; x_i x_{i+1}; 1 \leq i \leq n-1\} \cup \{x_{i,j} x_{i,j+1}; 1 \leq i \leq n; 1 \leq j \leq m-3\} \cup \{x_i x_{i,1}; 1 \leq i \leq n\} \cup \{x_1 x_{n,m-2}; x_{i+1} x_{i+1,1}; n \leq i \leq n\}$. Thus, the order and size of this graph are $p = |V(C_n \supseteq C_m)| = n + n(m-2), q = |E(C_n \supseteq C_m)| = mn$. Since all edges in C_n joint with one edge in C_m , it gives $\Delta(C_n \supseteq C_m) = 4$.

By Observation 1, $\chi(C_n \supseteq C_m) \geq \min\{r, \Delta(C_n \supseteq C_m)\} = \min\{r, 4\}$. To find the vertex r -dynamic chromatic number of $(C_n \supseteq C_m)$, we define three cases, namely for $\chi(C_n \supseteq C_m), \chi_d(C_n \supseteq C_m), \chi_3(C_n \supseteq C_m)$ and $\chi_4(C_n \supseteq C_m)$.

For $\chi(C_n \supseteq C_m), \chi_d(C_n \supseteq C_m)$, the lower bound $\chi_1(C_n \supseteq C_m) \geq \min\{2, 4\} = 2$. We will show that $\chi_1(C_n \supseteq C_m) \leq 3$, by defining a map $c_4 : V(C_n \supseteq C_m) \rightarrow \{1, 2, 3, \dots, k\}$ where $n \geq 3, m \geq 3$ by the following :

• For n and m even,

$$c_4(x_i) = \begin{cases} 1, & i \equiv 1 \pmod{2}, 1 \leq i \leq n, \\ 2, & i \equiv 0 \pmod{2}, 1 \leq i \leq n. \end{cases}$$

$$c_4(x_{i,j}) = \begin{cases} 2, & i \equiv 1 \pmod{2}, j \equiv 1 \pmod{2}, 1 \leq i \leq n, \\ & 1 \leq j \leq m-2, \\ 1, & i \equiv 1 \pmod{2}, j \equiv 0 \pmod{2}, 1 \leq i \leq n, \\ & 2 \leq j \leq m-2, \\ 2, & i \equiv 2 \pmod{2}, j \equiv 2 \pmod{2}, 1 \leq i \leq n, \\ & 1 \leq j \leq m-2, \\ 1, & i \equiv 0 \pmod{2}, j \equiv 1 \pmod{2}, 2 \leq i \leq n, \\ & 1 \leq j \leq m-2, \end{cases}$$

- For n and m odd,

$$c_4(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), 1 \leq i \leq n-1, \\ 2, & i \equiv 0(\text{mod } 2), 1 \leq i \leq n-1, \\ 3, & i = n. \end{cases}$$

$$c_4(x_{i,j}) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), j \equiv 0(\text{mod } 2), 1 \leq i \leq n-1, 2 \leq j \leq m-2, \\ 1, & i \equiv 0(\text{mod } 2), j \equiv 1(\text{mod } 2), 2 \leq i \leq n, 1 \leq j \leq m-3, x_{i-1}, \\ 2, & i \equiv 1(\text{mod } 2), j \equiv 1(\text{mod } 2), 1 \leq i \leq n, 1 \leq j \leq m-3, \\ 2, & i \equiv 0(\text{mod } 2), j \equiv 0(\text{mod } 2), 2 \leq i \leq n-1, 1 \leq j \leq m-3, \\ 3, & x_n; x_{i,m-3}; 1 \leq i \leq n-3 \end{cases}$$

- For $n = \text{odd}$ and $m = \text{even}$,

$$c_4(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), 1 \leq i \leq n-1, \\ 2, & i \equiv 0(\text{mod } 2), 1 \leq i \leq n-1, \\ 3, & i = n. \end{cases}$$

$$c_4(x_{i,j}) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), j \equiv 0(\text{mod } 2), 1 \leq i \leq n, 1 \leq j \leq m-2, \\ 2, & i \equiv 0(\text{mod } 2), j \equiv 1(\text{mod } 2), 1 \leq i \leq n. \end{cases}$$

- For $n = \text{even}$ and $m = \text{odd}$,

$$c_4(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), 1 \leq i \leq n, \\ 2, & i \equiv 0(\text{mod } 2), 1 \leq i \leq n. \end{cases}$$

$$c_4(x_{i,j}) = \begin{cases} 1, & i \equiv 1(\text{mod } 2), j \equiv 0(\text{mod } 2), 1 \leq i \leq n, 2 \leq j \leq m-3, \\ 1, & i \equiv 0(\text{mod } 2), j \equiv 1(\text{mod } 2), 1 \leq i \leq n, 1 \leq j \leq m-2, \\ 2, & i \equiv 1(\text{mod } 2), j \equiv 1(\text{mod } 2), \\ 2, & i \equiv 0(\text{mod } 2), j \equiv 0(\text{mod } 2), \\ 3, & 1 \leq i \leq n; j = m-2 \end{cases}$$

- For $\chi_{r=2}$, $n = 3$ and $m = 3, 6$,

$$c_4(x_i) = i; 1 \leq i \leq 3$$

$$c_4(x_{i,j}) = \begin{cases} 3, & i = 1, j = 1 \\ 1, & i = 2, j = 1 \\ 2, & i = 3, j = 1 \end{cases}$$

It easy to see that c_4 gives $\chi(C_n \supseteq C_m) \leq 3$ and $\chi_d(C_n \supseteq C_m) \leq 3$. Thus $\chi(C_n \supseteq C_m) = 3$ and $\chi_d(C_n \supseteq C_m) = 3$.

For $r = 3$, the lower bound $\chi_3(C_n \supseteq C_m) \geq \min\{3, 4\} = 3$. We will show that $\chi_3(C_n \supseteq C_m) \leq 4$, by defining a map $c_5 : V(C_n \supseteq C_m) \rightarrow \{1, 2, 3, \dots, k\}$ where $n \geq 3$ and $m \geq 3$ by the following:

- For $n = 3$ and $m = 3, 6$

$$c_5(x_i) = i$$

$$c_5(x_{i,j}) = \begin{cases} 3214; i = 1, \\ 1324; i = 2, \\ 2134; i = 3, \end{cases}$$

$$c_5(x_i) = i, c_5(x_{i,j}) = 4; n = 3, m = 3.$$

- For $n = 3k$ and $m = 3k$

$$c_5(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 3), \\ 2, & i \equiv 2(\text{mod } 3), \\ 3, & i \equiv 3(\text{mod } 3), \end{cases}$$

$$c_5(x_{i,j}) = \begin{cases} 321321 \dots 3214, & i \equiv 1(\text{mod } 3), \\ 132132 \dots 1324, & i \equiv 2(\text{mod } 3), \\ 213213 \dots 2134, & i \equiv 3(\text{mod } 3), \end{cases}$$

It is easy to that c_5 gives $\chi_3(C_n \supseteq C_m) \leq 4$. Thus $\chi_3(C_n \supseteq C_m) = 4$.

For $r = 4$, the lower bound $\chi_4(C_n \supseteq C_m) \geq \min\{4, 4\} = 4$. We will show that $\chi_4(C_n \supseteq C_m) \leq 5$, by defining a map $c_6 : V(C_n \supseteq C_m) \rightarrow \{1, 2, 3, \dots, k\}$ where $n \geq 3$ and $m \geq 3$ by the following:

- For $n = 3k$ and $m = 3k, 1 \leq i \leq n, 1 \leq j \leq m-2$

$$c_6(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 3), \\ 2, & i \equiv 2(\text{mod } 3), \\ 3, & i \equiv 3(\text{mod } 3). \end{cases}$$

$$c_6(x_{i,j}) = \begin{cases} 421421 \dots 4215, & i \equiv 1(\text{mod } 3), \\ 432432 \dots 4325, & i \equiv 2(\text{mod } 3), \\ 413413 \dots 4135, & i \equiv 3(\text{mod } 3), \end{cases}$$

- For $n = 3k$ and $m = 3k+1, 1 \leq i \leq n, 1 \leq j \leq m-2$

$$c_6(x_i) = \begin{cases} 1, & i \equiv 1(\text{mod } 3), \\ 2, & i \equiv 2(\text{mod } 3), \\ 3, & i \equiv 3(\text{mod } 3). \end{cases}$$

$$c_6(x_{i,j}) = \begin{cases} 421421 \dots 42145, & i \equiv 1(\text{mod } 3), \\ 432432 \dots 43245, & i \equiv 2(\text{mod } 3), \\ 415415 \dots 41545, & i \equiv 3(\text{mod } 3), \end{cases}$$

It is easy to that c_6 gives $\chi_4(C_n \supseteq C_m) \leq 5$. Thus $\chi_4(C_n \supseteq C_m) = 5$.

For $r = 5$, the lower bound $\chi_5(C_n \supseteq C_m) \geq \min\{5, 4\} = 4$. We will show that $\chi_5(C_n \supseteq C_m) \leq 5$, by defining a map $c_7 : V(C_n \supseteq C_m) \rightarrow \{1, 2, 3, \dots, k\}$ where $n \geq 2$ by the following:

- For $n = 5$ and $m = 3k$

$$c_7(x_i) = i$$

$$c_7(x_{i,j}) = \begin{cases} 421421 \dots 421, 5 & i = 1, \\ 432432 \dots 432, 1 & i = 2, \\ 543543 \dots 543, 1 & i = 3, \\ 254254 \dots 254, 3 & i = 4, \\ 215215 \dots 215, 3 & i = 5, \end{cases}$$

- For $n = 5$ and $m = 3k + 1$

$$c_7(x_i) = i$$

$$c_7(x_{i,j}) = \begin{cases} 421\ 421 \cdots 421, 5 & i = 1, \\ 324\ 324 \cdots 324, 1 & i = 2, \\ 543\ 543 \cdots 543, 1 & i = 3, \\ 254\ 254 \cdots 254, 3 & i = 4, \\ 215\ 215 \cdots 215, 3 & i = 5, \end{cases}$$

It is easy that c_7 gives $\chi_5(C_n \triangleright C_m) \leq 5$. Thus $\chi_5(C_n \triangleright C_m) = 5$. Since for $r \geq 5$, we have $r \geq \Delta(C_n \triangleright C_m)$. By Observation 1, $\chi_r(C_n \triangleright C_m) = \chi_5(C_n \triangleright C_m) = 5$. It concludes the proof. \square

Theorem 4. Let G be a comb product denote by $C_n \triangleright C_m$ for $n \geq 3$ and $m \geq 3$, edges r -dynamic chromatic number of $(C_n \triangleright C_m)$ is :

$$\chi_{1 \leq r \leq 3}(C_n \triangleright C_m) = 4$$

Proof. The graph $(C_n \triangleright C_m)$ is a connected graph with vertex set $V(C_n \triangleright C_m) = \{x_i, ; 1 \leq i \leq n\} \cup \{x_{i,j}; 1 \leq i \leq n; 1 \leq j \leq m - 2\}$ and edge set $E(C_n \triangleright C_m) = \{x_n x_1; x_i x_{i+1}; 1 \leq i \leq n - 1\} \cup \{x_{i,j} x_{i,j+1}; 1 \leq i \leq n; 1 \leq j \leq m - 3\} \cup \{x_i x_{i,1}; 1 \leq i \leq n\} \cup \{x_1 x_{n,m-2}; x_{i+1} x_{i+1,1}; n \leq i \leq n\}$. Thus, the order and size of this graph are $p = |V(C_n \triangleright C_m)| = n + n(m - 2)$, $q = |E(C_n \triangleright C_m)| = mn$. Since all edges in C_n joint with one edge in C_m , it gives $\Delta(C_n \triangleright C_m) = 4$.

By Observation 1, $\chi_r(C_n \triangleright C_m) \geq \min\{\chi(C_n \triangleright C_m), r\} + 1 = \min\{r, d(u) + d(v) - 2\}$. To find the exact value of r -dynamic chromatic number of $(C_n \triangleright C_m)$, we define three cases, namely for $\chi_{1 \leq r \leq 4}(C_n \triangleright C_m)$, $\chi_5(C_n \triangleright C_m)$.

For $r = 1$, the lower bound $\Delta(G) \leq \chi(G) \leq \Delta(G) + 1$, that $\chi(C_n \triangleright C_m) \geq 3$. Furthermore, to show that $\chi(C_n \triangleright C_m) \leq 3$ with coloring edges $E(C_n \triangleright C_m)$ as in function c_8 . Let $D = \{1, 2, \dots, k\}$ is set of color from k -coloring that c_1 the function by defining a map edges coloring D , $c_8 : E(C_n \triangleright C_m) \rightarrow D$, so mapping each edges to set color D , by the following that:

$$c_8(x_i x_{i+1}) = \begin{cases} 1, & i \text{ odd } 1 \leq i \leq n - 1 \\ 2, & i \text{ even } 1 \leq i \leq n - 1 \end{cases}$$

$$c_8(x_n x_{i+1}) = \begin{cases} 2, & i \text{ odd } 1 \leq i \leq n - 1 \\ 3, & i \text{ even } 1 \leq i \leq n - 1 \end{cases}$$

$$c_8(x_i y_{1,j}) = 4, 1 \leq i \leq n \ 1 \leq j \leq m$$

- For $n = \text{even}$

$$c_8(y_{i,j} y_{i,j+1}) = \begin{cases} 1, & j \equiv 1(\text{mod } 3), \\ & 1 \leq i \leq m - 3 \\ \vdots \\ 2, & j \equiv 2(\text{mod } 3), \\ & 1 \leq i \leq m - 3 \\ \vdots \\ 4, & j \equiv 0(\text{mod } 3), \\ & 1 \leq i \leq m - 3 \end{cases}$$

$$c_8(x_{i+1} y_{i,m-2}) = 3$$

$$c_8(x_1 y_{n,m-2}) = 3$$

- For $n = \text{even}$

$$c_8(y_{i,j} y_{i,j+1}) = \begin{cases} 1, & j \equiv 1(\text{mod } 3), i = n \\ & 1 \leq i \leq n - 2; \\ \vdots \\ & 1 \leq i \leq m, \\ \vdots \\ 2, & j \equiv 2(\text{mod } 3), \\ & i = n - 1 \\ \vdots \\ & 2 \leq i \leq n - 2; \\ & 2 \leq i \leq m, \\ \vdots \\ 3, & j \equiv 2(\text{mod } 3), \\ & n - 1 \leq i \leq n \\ \vdots \\ 4, & 1 \leq i \leq n; \\ \vdots \\ & 1 \leq j \leq m \end{cases}$$

$$c_8(x_{i+1} y_{i,m-2}) = 3, 1 \leq i \leq n - 2$$

$$c_8(x_1 y_{n,m-2}) = 2$$

$$c_8(x_{n-1} y_{n-1,m-2}) = 1$$

From function coloring c_8 seen that the chromatic number is $\chi(C_n \triangleright C_m) \leq 4$. Because $\chi(C_n \triangleright C_m) \leq 4$ and $\chi(C_n \triangleright C_m) \geq 4$, then $\chi(C_n \triangleright C_m) = 4$, so that $\chi(C_n \triangleright C_m) = \chi_2(C_n \triangleright C_m) = \chi_3(C_n \triangleright C_m) = 4$.

CONCLUSIONS

We have found some edge and vertex r -dynamic chromatic number of several graphs, namely comb product of graph $C_n \triangleright P_2$ and $C_n \triangleright C_m$. It is interesting to characterize a property of any graph operation to have an exact value or upper bound of their r -dynamic chromatic numbers.

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